

# The preserving non-falsity companion of the Nilpotent Minimum Logic

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We recall that a logic  $L$  is said to be *paraconsistent* with respect to a negation connective  $\neg$  when it contains a  $\neg$ -contradictory but not trivial theory. Assuming that  $L$  is (at least) Tarskian, this is equivalent to say that the  $\neg$ -explosion rule  $\frac{\varphi \quad \neg\varphi}{\psi}$  is not valid in  $L$ .

The 3-valued logic  $J_3$  introduced by D'Ottaviano and da Costa in [2] is one of the well known paraconsistent logics and it can be defined (up to language) as the logic given by the matrix  $\langle \mathbf{MV}_3, \{\frac{1}{2}, 1\} \rangle$  where  $\mathbf{MV}_3$  is the 3 element MV-chain. Notice that  $J_3$  is strongly related with the 3-valued Łukasiewicz logic  $L_3$  as  $\langle \mathbf{MV}_3, \{1\} \rangle$  is a matrix semantics for  $L_3$ . Moreover, these two logics are equivalent deductive systems in the Blok-Pigozzy sense [1]. Notice that, while  $L_3$  is explosive and truth-preserving (1 being full truth),  $J_3$  is paraconsistent and non-falsity-preserving, because it preserves every element different from 0 (0 being false). We call  $J_3$  the *non-falsity companion* of  $L_3$ .

The *nilpotent minimum logic*, NML for short, was firstly introduced by Esteva and Godo in [3] in order to formalize the logic of the nilpotent minimum t-norm, that was defined by Fodor in [4] as an example of an involutive left continuous t-norm which is not continuous. NML is obtained from the monoidal t-norm logic MTL defined in [3], by adding the involutive condition axiom (INV)  $\neg\neg\varphi \rightarrow \varphi$  and the (weak) nilpotent minimum condition axiom (WNM)  $(\psi * \varphi \rightarrow \perp) \vee (\psi \wedge \varphi \rightarrow \psi * \varphi)$ . It is well known that NML is algebraizable and the class  $\mathbf{NML}$  of all nilpotent minimum algebras is its equivalent algebraic quasivariety semantics [3]. Moreover, NML is sound and strong complete with respect the standard NM-algebra  $[0, 1]_{\mathbf{NM}}$  [7]. That is, NML is the logic defined by the matrix  $\langle [0, 1]_{\mathbf{NM}}, \{1\} \rangle$ . The aim of this talk is to axiomatize and characterize the non-falsity companions of NML and its axiomatic extensions.

Let  $\mathbf{A}$  be a subalgebra of  $[0, 1]_{\mathbf{NM}}$ , then the finitary logic  $L$  defined by  $\langle \mathbf{A}, \{1\} \rangle$  is an axiomatic extension (not necessarily proper) of  $NML$ . We call  $\text{nf-}L$  the non-falsity companion of  $L$ . That is,  $\text{nf-}L$  is the finitary logic defined by the matrix  $\langle \mathbf{A}, (0, 1] \cap A \rangle$ . Consider now the following *restricted* inference rule, which is intended for axiomatising  $\text{nf-}L$ :

- Restricted Square Modus Ponens for  $L$  (r-MP<sup>2</sup> for  $L$ ):

From  $\varphi$  and  $\varphi \rightarrow \neg(\neg\psi)^2$  derive  $\psi$ , whenever  $\vdash_L \varphi \rightarrow \neg(\neg\psi)^2$ .

It is not hard to see that from (r-MP<sup>2</sup> for  $L$ ) we can derive the following restricted version of Modus Ponens:

- Restricted Modus Ponens for  $L$  (r-MP for  $L$ ):

From  $\varphi$  and  $\varphi \rightarrow \psi$  derive  $\psi$ , whenever  $\vdash_L \varphi \rightarrow \psi$

Note that both inference rules involve conditions on the derivability of formulas in the logic  $L$ . Since any axiomatic extension of NML is complete w.r.t at most two subalgebras of  $[0, 1]_{\mathbf{NM}}$  [5] we obtain the following result.

**Theorem 1.** *Let  $L$  be an axiomatic extension of NML. The following axiomatization*

- *Axioms: those of  $L$*
- *Rules: Adjunction  $\frac{\varphi \quad \psi}{\varphi \wedge \psi}$  and  $(r\text{-MP}^2)$  for  $L$*

*is a sound and complete axiomatisation of  $\text{nf-}L$ .*

For the case of finite-valued axiomatic extensions  $\text{NM}_n$ , unlike the Łukasiewicz case [1, Th.5.2], we prove that  $\text{nf-NM}_n$  is not equivalent to  $\text{NM}_n$ . With an abuse of language,  $\mathcal{N}_k$  denotes the matrix  $\langle \text{NM}_k, \{1\} \rangle$  and  $\mathcal{J}_k$  will denote the matrix  $\langle \text{NM}_k, \{\frac{1}{k-1}, \frac{2}{k-1}, \dots, 1\} \rangle$  where  $\text{NM}_k$  is the  $k$ -element NM-chain. It is shown in [6] that any finitary extension of  $\text{NM}_n$  is complete w.r.t. following set of matrices  $\{\mathcal{N}_{2k}, \mathcal{N}_{2m+1}, \mathcal{N}_2 \times \mathcal{N}_{2r+1}\}$  for some  $0 \leq m \leq r \leq k \leq n$ . For the case of  $\text{nf-NM}_n$  we cannot accomplish this reduction, but the following one that is restricted to finitary extensions defined by finite products of  $\mathcal{J}_k$ 's.

**Theorem 2.** *Let  $L$  be a finitary extension of  $\text{nf-NML}$  defined by  $\mathcal{J}_{k_1} \times \dots \times \mathcal{J}_{k_s}$ . Then  $L$  is complete w.r.t a finite set of the following matrices:*

- (i)  $\mathcal{J}_n$  for some positive integer  $n > 1$ .
- (ii)  $\mathcal{J}_n \times \mathcal{J}_k$  for some positive integers  $n \neq k$ .
- (iii)  $\mathcal{J}_{2n} \times \mathcal{J}_{2k} \times \mathcal{J}_{2l+1}$  for some positive integers  $l < n < k$ .
- (iv)  $\mathcal{J}_{2n} \times \mathcal{J}_{2m+1} \times \mathcal{J}_{2l+1}$  for some positive integers  $m < n$  and  $m < l$ .

Moreover every different matrix of these four types defines a different logic

Finally, next result characterizes all finite maximal paraconsistent extensions  $\text{nf-NML}$

**Theorem 3.** *The only finite matrices defining maximal paraconsistent extensions of  $\text{nf-NML}$  are  $\mathcal{J}_3$ ,  $\mathcal{J}_4$  and  $\mathcal{J}_3 \times \mathcal{J}_4$ .*

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